

SDP

$$\min_x \langle C, X \rangle$$

$$\text{s.t. } X \succeq_{\text{SD}} 0$$

$$A(X) = b$$

$$X \in S_{\text{H}}^n$$

$$b - A(X) \in S_{\text{H}}^m$$

$$J(X, \lambda, \nu) = \langle C, X \rangle - \langle \lambda, X \rangle - \langle \nu, b - A(X) \rangle$$

$$= \langle C - \lambda + A^T(\nu), X \rangle - \langle b, \nu \rangle$$

$$= \begin{cases} -\infty & C + A^T(\nu) - \lambda \neq 0 \\ -\langle b, \nu \rangle & C + A^T(\nu) - \lambda = 0 \end{cases}$$

$$\max_{\nu} -\langle b, \nu \rangle$$

$$\text{s.t. } C + A^T(\nu) \succeq_{\text{SD}} 0$$

$$\max_{\nu} -\langle b, \nu \rangle$$

$$\text{s.t. } C + A^T(\nu) \in (S_{\text{H}}^n)^* = S_{\text{H}}^n$$

Weak duality fin value of P $\geq$ value of D

if the primal problem is finite-feasible and dual problem is feasible then the finite value of P is bounded by value of D

let $\{X_k\}$ be a feasible sequence ($X_k \in S_{\text{H}}^n$ $\lim_{k \rightarrow \infty} A(X_k) = b$)

$$\liminf_{k \rightarrow \infty} \langle C, X_k \rangle - \langle b, \nu \rangle = \liminf_{k \rightarrow \infty} \langle C, X_k \rangle + \lim_{k \rightarrow \infty} \langle A(X_k), \nu \rangle \geq \liminf_{k \rightarrow \infty} \underbrace{\langle C + A^T(\nu), X_k \rangle}_{\substack{S_{\text{H}}^n \\ \succeq 0}} \geq 0$$

Regular Duality fin value of P = value of D

The dual problem is feasible has finite value d^*

The primal problem is finite-feasible, has finite val p^*

$$\left. \begin{array}{l} \text{The dual problem is feasible has finite value } d^* \\ \text{The primal problem is finite-feasible, has finite val } p^* \end{array} \right\} p^* = d^*$$

$p^* \downarrow$

$$A^T(\nu) + zC \in S_{\text{H}}^n$$

$$\downarrow$$

$$-\langle b, \nu \rangle \leq z d^*$$

$$\text{for } z > 0$$

if D is feasible and optimal value = d^*

then $\forall \nu$ s.t. $C + A^T(\nu) \in S_{\text{H}}^n$, we have $-\langle b, \nu \rangle \leq d^*$

then $\forall \nu \succeq 0$ s.t. $C + A^T(\nu/2) \in S_{\text{H}}^n$, we have $-\langle b, \nu/2 \rangle \leq d^*$

then $\forall \nu \succeq 0$ s.t. $2C + A^T(\nu) \in S_{\text{H}}^n$, we have $-\langle b, \nu \rangle \leq 2d^*$

$$zC + A^T(\nu) \in S_{\text{H}}^n$$

$$\downarrow$$

$$-\langle b, \nu \rangle \leq z d^*$$

$$\text{for } z > 0$$

also $\forall \nu$ s.t. $A^T(\nu) \in S_{\text{H}}^n$, we have $-\langle b, \nu \rangle \leq 0$

if $\exists \tilde{\nu}$ s.t. $A^T(\tilde{\nu}) \in S_{\text{H}}^n$ $-\langle b, \tilde{\nu} \rangle > 0$

$\forall \nu$ feasible for D

$$A^T(\nu + \tilde{\nu}) \in S_{\text{H}}^n$$

$$-\langle b, \nu + \tilde{\nu} \rangle \rightarrow +\infty$$

con-contradiction d^* is optimal value

$$\left\{ \begin{array}{l} A^T(v) + zC \in S^+ \\ z > 0 \end{array} \right\} \Rightarrow \left\{ \langle v, b \rangle \leq z d^* \right\}$$

$\therefore$ the system $\left\{ \begin{array}{l} z > 0 \\ A^T(v) + zC \in S^+ \\ -\langle v, b \rangle > z d^* \end{array} \right\}$ is infeasible

$\therefore$ the system $\left\{ \begin{array}{l} \begin{bmatrix} A^T & C \\ 0 & 1 \end{bmatrix} \begin{bmatrix} v \\ z \end{bmatrix} \in \begin{bmatrix} S^+ \\ \mathbb{R}_+ \end{bmatrix} \\ \langle [v, z], [b, d^*] \rangle < 0 \end{array} \right\}$ is infeasible

by Farkas Lemma

$\therefore$ the system $\left\{ \begin{array}{l} \begin{bmatrix} A & 0 \\ C^T & 1 \end{bmatrix} \begin{bmatrix} x \\ z \end{bmatrix} = \begin{bmatrix} b \\ d^* \end{bmatrix} \\ \begin{bmatrix} x \\ z \end{bmatrix} \in \begin{bmatrix} S^+ \\ \mathbb{R}_+ \end{bmatrix} \end{array} \right\}$ is infeasible

$\therefore \exists \{x_k\} \{z_k\}$ s.t. $\underbrace{\lim_{k \rightarrow \infty} A(x_k) = b}_{\text{in-feas}} \quad \underbrace{\lim_{k \rightarrow \infty} \langle C, x_k \rangle + z_k = d^*}_{\substack{\text{finite value} \leq d^* \\ \liminf \langle C, x \rangle \leq d^*}}$

by weak duality, $\nexists$ feasible sequence $\{x_k\}$. $\liminf_{k \rightarrow \infty} \langle C, x \rangle > d^*$

$\therefore \liminf_{k \rightarrow \infty} \langle C, x \rangle = d^*$

$\Rightarrow \uparrow$

if P is infeasible and the finite value is p^*

$\exists \{x_k\}$ s.t. $x_k \in S^+ \quad \lim_{k \rightarrow \infty} A(x_k) = b \quad \lim_{k \rightarrow \infty} \langle C, x \rangle = p^*$

assume D is infeasible (we'll get a contradiction)

if D is infeasible, we have

$$A^T(v) + zC \in S^+ \Rightarrow z \leq 0$$

for any $A^T(v) + zC \in S^+ \quad z > 0$, v/z is feasible for D

$\therefore$ the system $\left\{ \begin{array}{l} A^T(v) + zC \in S^+ \\ z > 0 \end{array} \right\}$ is infeasible

$\therefore$ the system $\left\{ \begin{array}{l} \begin{bmatrix} A^T & C \\ 0 & 1 \end{bmatrix} \begin{bmatrix} v \\ z \end{bmatrix} \in \begin{bmatrix} S^+ \\ \mathbb{R}_+ \end{bmatrix} \\ \langle [v, z], [0, -1] \rangle < 0 \end{array} \right\}$ is infeasible

by Farkas lemma

then the system $\left\{ \begin{array}{l} \begin{bmatrix} A \\ C^T \end{bmatrix} \begin{bmatrix} x \\ z \end{bmatrix} = \begin{bmatrix} b \\ -1 \end{bmatrix} \\ x \in S^+ \end{array} \right\}$ is infeasible

$\therefore \exists \{x_k\}$ s.t. $\lim_{k \rightarrow \infty} A(x_k) = b \quad \lim_{k \rightarrow \infty} \langle C, x \rangle = -1$

then consider the sequence $\{x_k + z_k x_k\}$ is a feasible sequence

and the value is unbounded below

$\therefore$ when P is lim-fea, D must be for, and $p^* = d^*$ when D is fea.

Strong Duality

if P is feasible, has a finite p^* and an interior point
then D is feasible, has $d^* = p^*$

$$\left. \begin{array}{l} P \text{ fea} + \text{int. point} \Rightarrow \lim p^* = p^* \\ P \text{ is lim-fea} \Rightarrow D \text{ is for, } d^* = \lim p^* \end{array} \right\} p^* = d^*$$

Strong Duality May Not Hold

$$\left. \begin{array}{l} \min_x z/x_2 \\ \text{s.t. } \begin{bmatrix} 0 & x_2 \\ x_1 & x_2 \end{bmatrix} \succeq 0 \end{array} \right\} \text{primal} \quad \left. \begin{array}{l} \min_x \langle \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}, x \rangle \\ \text{s.t. } \langle \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}, x \rangle = 0 \\ x \succeq_{\text{SD}} 0 \end{array} \right\} p^* = 0 \text{ when } x_2 = x_1$$

$$\swarrow \text{dual} \quad \left. \begin{array}{l} \max_{\text{DU}} 0 \\ \text{s.t. } \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} + v \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} \succeq_{\text{SD}} 0 \end{array} \right\} d \text{ is infeasible}$$

$$P \text{ is limit feasible with } x_k = \begin{bmatrix} 1/k & 1/k \\ 1/k & 1/k \end{bmatrix} \succeq_{\text{SD}} 0$$

$$\lim_{k \rightarrow \infty} \langle \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}, x \rangle = \lim_{k \rightarrow \infty} \frac{1}{k} = 0$$

$$\text{the limit value} = \liminf_{k \rightarrow \infty} z_k = -\infty$$

$$\left(\begin{array}{l} \text{if } P \text{ has no interior, } \lim p^* \neq p^* \\ \text{if } P \text{ is lim-fea, but the lim value is not finite} \end{array} \right)$$

optimality Conditions

if P is for and has an interior-point

$$\begin{aligned} \langle C, x^* \rangle &= -\langle b, v^* \rangle = \inf_{x \in S^*} [\langle C, x \rangle - \langle \lambda^*, x \rangle - \langle v^*, b - A(x) \rangle] \\ f(x^*) &= g(v^*, \lambda^*) \leq \langle C, x^* \rangle - \langle \lambda^*, x^* \rangle - \langle v^*, b - A(x^*) \rangle \\ &\leq \langle C, x^* \rangle \end{aligned}$$

$$\lambda^* = C + A^T(v^*)$$

optimality condition

primal fea : $x \succeq_{SD} 0 \quad A(x) = b$

Dual fea : $\lambda \succeq_{SD} 0$

Complementarity : $X\lambda = 0$

gradient vanish : $\lambda = C + A^T(v)$

$$X\lambda = 0 \Leftrightarrow \langle X, \lambda \rangle = 0$$

$$(X \in S^+, \lambda \in S^+)$$

$$1' \quad \text{if } \langle X, \lambda \rangle = 0$$

$$X = PP^T, \lambda = QQ^T$$

$$\langle X, \lambda \rangle = \text{tr}(X\lambda) = \text{tr}(PP^TQQ^T)$$

$$= \text{tr}(P^TQQP)$$

$$= \langle P^TQ, P^TQ \rangle$$

$$\geq 0$$

$$\therefore P^TQ = 0$$

$$X\lambda = PP^TQQ^T = P0Q^T = 0$$

$$2' \quad \text{if } X\lambda = 0$$

$$\text{then } U^T X \lambda V = 0 \quad \forall U, V$$

$$\langle U, U^T X \lambda V \rangle = 0 \quad \forall U, V$$

$$\therefore X\lambda = 0$$

Log-Barrier

$$\min_x \langle C, x \rangle$$

$$\text{s.t. } A(x) = b$$

$$x \succ 0$$

↓

$$\min_x \langle C, x \rangle - \frac{1}{\epsilon} \log \det(x)$$

$$\text{s.t. } A(x) = b$$

$$\mathcal{L}(x, \lambda, v) = \langle C, x \rangle - \langle \lambda, x \rangle - \langle v, b - A(x) \rangle$$

$$\Rightarrow x^*(t)$$

$$\mathcal{L}(x, \tilde{v}) = t \langle C, \tilde{x} \rangle - \log \det(\tilde{x}) - \langle v, b - A(\tilde{x}) \rangle$$

$$= \langle \tilde{x}, tC + A^T(v) \rangle - \log \det(\tilde{x}) - \langle v, b \rangle$$

$$\text{optimality condition: } \begin{cases} \nabla \mathcal{L} = tC + A^T(v) - \tilde{x}^{-1} = 0 \\ A(\tilde{x}) = b \quad \tilde{x} \succ 0 \end{cases}$$

$$C + A^T(\tilde{v}/t) - \frac{1}{t} \tilde{x}^{-1} = 0$$

$x^*(t)$ minimize the Lagrangian

$$\mathcal{L}(x, \lambda^*(t), v^*(t)) = \langle C, x \rangle - \langle \lambda^*(t), x \rangle - \langle v^*(t), b - A(x) \rangle$$

$$\text{where } \lambda^*(t) = \frac{1}{t} x^*(t)^{-1} \quad v^*(t) = \tilde{v}/t$$

$$p^* \geq g(\lambda^*(t), v^*(t))$$

$$= \min_x \mathcal{L}(x, \lambda^*(t), v^*(t))$$

$$= \mathcal{L}(x^*(t), \lambda^*(t), v^*(t))$$

$$= \langle C, x^*(t) \rangle - \langle \frac{1}{t} x^*(t)^{-1}, x^*(t) \rangle - \langle v^*(t), b - A(x^*(t)) \rangle$$

$$= \langle C, x^*(t) \rangle - \underbrace{n/t}_{\text{duality gap}}$$

$$X = X^*(t), \quad \Lambda = \Lambda^*(t) \quad V = V^*(t) \text{ satisfies}$$

$$\text{primal: } X \geq 0 \quad A(X) = b$$

$$\text{dual: } \Lambda \geq 0$$

$$\text{approximate complementarity: } X\Lambda = \frac{1}{t}I$$

$$\text{gradient vanishes: } C - \Lambda + A^T(V) = 0$$

• Search Direction

$$r_t(X, \Lambda, V) = \left[\begin{array}{c} \langle A_1, X \rangle - b_1 \\ \vdots \\ \langle A_m, X \rangle - b_m \\ X\Lambda - \frac{1}{t}I \\ C - \Lambda + A^T(V) \end{array} \right] \left. \begin{array}{l} \text{primal residual} \\ \\ \text{centrality residual} \\ \text{dual residual} \end{array} \right\}$$

$$r_t(X + \delta X, \Lambda + \delta \Lambda, V + \delta V) \approx r_t(X, \Lambda, V) + Dr_t(X, \Lambda, V) \begin{bmatrix} \delta X \\ \delta \Lambda \\ \delta V \end{bmatrix} := 0$$

• Path-Following Method

while true:

$$\text{if } C - \Lambda + A^T(V) = 0 \text{ and } A(X) = b \text{ and } \langle \Lambda, X \rangle \leq \epsilon \mu$$

return X, Λ, V

$$t = m / \langle \Lambda, X \rangle \cdot \mu$$

solve the system (log-barrier)

$$\begin{cases} C - (\Lambda + \delta \Lambda) + A^T(V + \delta V) = 0 \\ A(X + \delta X) = b \\ (X + \delta X)(\Lambda + \delta \Lambda) = \frac{1}{t}I \end{cases}$$

line search on residual norm to get S

$$X \leftarrow X + S \cdot \delta X$$

$$\Lambda \leftarrow \Lambda + S \cdot \delta \Lambda$$

$$V \leftarrow V + S \cdot \delta V$$