



Semidefinite programming

$$\begin{aligned} \min & \langle C, X \rangle \\ \text{st.} & A(X) = b \\ & X \succeq_{SD} 0 \end{aligned}$$

$$\begin{aligned} \mathcal{L}(X, \lambda, V) &= \langle C, X \rangle - \langle \lambda, X \rangle - \langle V, A(X) - b \rangle \\ &= \langle C - \lambda - A^T(V), X \rangle + \langle b, V \rangle \\ &= \begin{cases} -\infty & C - \lambda - A^T(V) \neq 0 \\ \langle b, V \rangle & C - \lambda - A^T(V) = 0 \end{cases} \end{aligned}$$

$$\begin{aligned} \max_{\lambda, V} & \langle b, V \rangle \\ \text{st.} & C - \lambda - A^T(V) = 0 \\ & \lambda \succeq_{SD} 0 \end{aligned}$$

Weak duality: if X is primal feasible and λ, V are dual feasible

$$\begin{aligned} \langle C, X \rangle &\geq \langle b, V \rangle \\ \langle C, X \rangle - \langle b, V \rangle &= \langle C, X \rangle - \langle A(X), V \rangle = \langle C - A^T(V), X \rangle = \langle \lambda, X \rangle \geq 0 \end{aligned}$$

$$\begin{aligned} p^* = g(\lambda^*, V^*) &= \inf_X \mathcal{L}(X, \lambda^*, V^*) \\ &= \inf_X \langle C, X \rangle - \langle \lambda^*, X \rangle - \langle A(X), V^* \rangle \\ &\leq \langle C, X^* \rangle - \langle \lambda^*, X^* \rangle - \langle A(X^*), V^* \rangle \\ &\leq p^* \end{aligned}$$

$$\begin{aligned} \langle A(X), V \rangle &= \sum_i v_i \lambda_i X \\ &= \langle \sum_i v_i \lambda_i, X \rangle \end{aligned}$$

eg. Strong duality of SDP may not hold

$$\begin{aligned} \min & \lambda_0 \\ \text{st.} & \begin{bmatrix} 0 & \lambda_0 \\ \lambda_0 & \lambda_0 \end{bmatrix} \succeq_{SD} 0 \end{aligned} \xrightarrow{\text{primal}} \begin{aligned} \min_X & \langle \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}, X \rangle \\ \text{st.} & \langle \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}, X \rangle = 0 \\ & X \succeq_{SD} 0 \end{aligned}$$

$\lambda_0 > 0 \implies 0 / \lambda_0 - \lambda_0 > 0$
the primal optimal is 0 (when $\lambda_0 = 0$)

primal is not strictly feasible

$$\begin{aligned} \max_{\lambda, V} & b \cdot V \\ \text{st.} & \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} - \lambda = V \cdot \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} \\ & \lambda \succeq_{SD} 0 \end{aligned} \left. \begin{array}{l} \\ \\ \end{array} \right\} \begin{aligned} & \begin{bmatrix} -V & 1 \\ 1 & 0 \end{bmatrix} \succeq_{SD} 0 \\ & -V > 0 \quad -V \cdot 0 - 1 > 0 \end{aligned} \text{ dual is infeasible} \end{aligned}$$

$$\begin{aligned} \min & \lambda_0 \\ \text{st.} & \begin{bmatrix} \lambda_0 & 1 \\ 1 & \lambda_0 \end{bmatrix} \succeq_{SD} 0 \end{aligned} \xrightarrow{\text{primal}} \begin{aligned} \min_X & \langle \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}, X \rangle \\ \text{st.} & \langle \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}, X \rangle = 2 \\ & X \succeq_{SD} 0 \end{aligned}$$

$\lambda_0 > 0 \implies \lambda_0 > 0$
 $\lambda_0 / \lambda_0 > 1$
the optimal value is 0 but not attained

$$\begin{aligned} \max & \lambda \cdot V \\ \text{st.} & \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} - \begin{bmatrix} 0 & V \\ V & 0 \end{bmatrix} = \lambda \\ & \lambda > 0 \end{aligned} \left. \begin{array}{l} \\ \\ \end{array} \right\} \begin{aligned} & \begin{bmatrix} 1 & -V \\ -V & 0 \end{bmatrix} \succeq 0 \\ & 0 = V > 0 \text{ only if } V > 0 \end{aligned} \text{ the dual is not strictly feasible} \\ & \text{primal optimal is 0, but attained} \end{aligned}$$

optimality condition

primal feasible: $X \succ_{\text{str}} 0$ $A(X) = b$

dual feasible: $\Lambda \succ_{\text{str}} 0$

complementary slackness: $X\Lambda = 0$

gradient vanishes: $A^T(U) + \Lambda = C$

$$\langle \Lambda, X \rangle = 0$$

$$\Lambda = PP^T \quad X = \alpha\alpha^T$$

$$0 = \text{tr}(\Lambda X) = \text{tr}(PP^T \alpha\alpha^T) = \text{tr}(P^T \alpha\alpha^T P) = \langle P^T \alpha, P\alpha \rangle$$

$$\therefore P^T \alpha = 0$$

$$\therefore \Lambda X = PP^T \alpha\alpha^T = P \alpha\alpha^T P^T = 0$$

Barrier of Primal and Dual

$$\min_x \langle C, X \rangle$$

$$\text{s.t. } A(X) = b$$

$$X \succ_{\text{str}} 0$$

↓

$$\min_x \langle C, X \rangle - \frac{1}{\epsilon} \log \det(X)$$

$$\text{s.t. } A(X) = b$$

$$\max_{\Lambda, U} \langle b, U \rangle$$

$$\text{s.t. } A^T(U) + \Lambda = C$$

$$\Lambda \succ_{\text{str}} 0$$

↓

$$\min_{\Lambda} -\langle b, U \rangle - \frac{1}{\epsilon} \log \det(\Lambda)$$

$$\text{s.t. } A^T(U) + \Lambda = C$$

$$\min_x \langle C, X \rangle - \frac{1}{\epsilon} \log \det X$$

$$\text{s.t. } A(X) = b$$

} primal barrier

$$\begin{aligned} \mathcal{L}(X, U) &= \langle C, X \rangle - \frac{1}{\epsilon} \log \det(X) - \langle U, A(X) - b \rangle \\ &= \langle C - A^T(U), X \rangle - \frac{1}{\epsilon} \log \det(X) + \langle U, b \rangle \end{aligned}$$

$$\nabla_X \mathcal{L} = C - \frac{1}{\epsilon} X^{-1} - A^T(U) = 0 \quad X = \frac{1}{\epsilon} (C - A^T(U))^{-1}$$

$$g(U) = \inf_X \mathcal{L}(X, U)$$

$$= \begin{cases} -\infty & \text{if } C - A^T(U) \not\succeq_{\text{str}} 0 \\ \frac{1}{\epsilon} \log \det \left(\frac{1}{\epsilon} (C - A^T(U))^{-1} \right) + \langle U, b \rangle & \text{otherwise} \end{cases}$$

if $C - A^T(U)$ has a negative eigenvalue λ_i , with eigenvector u

let $X = \alpha \cdot uu^T$

$$\mathcal{L}(X, U) = \langle C - A^T(U), \alpha uu^T \rangle - \frac{1}{\epsilon} \log \det \alpha uu^T + \langle U, b \rangle \rightarrow -\infty \text{ as } \alpha \rightarrow +\infty$$

$$\max_{\Lambda, U} \frac{1}{\epsilon} \log \det \left(\frac{1}{\epsilon} (C - A^T(U))^{-1} \right) + \langle U, b \rangle$$

$$\text{s.t. } C - A^T(U) \succ_{\text{str}} 0$$

$$\Rightarrow \left. \begin{aligned} \max_{\Lambda, U} & \frac{1}{\epsilon} \log \det(\Lambda) + \langle U, b \rangle \\ \text{s.t. } & C - A^T(U) = \Lambda \\ & \Lambda \succ_{\text{str}} 0 \end{aligned} \right\} \text{dual barrier}$$

KKT for primal barrier

primal feasible: $A(X) = b \quad X \succ_{\text{str}} 0$

dual feasible: $\Lambda + A^T(U) = C \quad \Lambda \succ_{\text{str}} 0$

gradient vanishes: $C - A^T(U) = \frac{1}{\epsilon} X^{-1} \Rightarrow X\Lambda = \frac{1}{\epsilon} I$

$$\begin{array}{l} \min_{\lambda} \quad -\langle b, v \rangle - \frac{1}{\epsilon} \log \det(\lambda) \\ \text{s.t.} \quad A^T(v) + \lambda = c \end{array} \quad \left. \vphantom{\begin{array}{l} \min_{\lambda} \\ \text{s.t.} \end{array}} \right\} \text{dual barrier}$$

$$\begin{aligned} \mathcal{L}(\lambda, v, x) &= -\langle b, v \rangle - \frac{1}{\epsilon} \log \det(\lambda) + \langle x, A^T(v) + \lambda - c \rangle \\ &= -\langle A(v) - b, v \rangle - \frac{1}{\epsilon} \log \det(\lambda) + \langle \lambda, x \rangle - \langle c, x \rangle \end{aligned}$$

$$\nabla_{\lambda} \mathcal{L} = -\frac{1}{\epsilon} \lambda^{-1} + x = 0 \quad \lambda^{-1} = \frac{1}{\epsilon} x$$

$$\begin{aligned} g(x) &= \min_{\lambda, v} \mathcal{L} \\ &= \begin{cases} -\langle b, v \rangle & A(v) - b = 0 \\ -\infty & \text{otherwise} \\ -\frac{1}{\epsilon} \log \det(\lambda) + \eta \frac{1}{\epsilon} \langle -c, x \rangle \end{cases} \end{aligned}$$

$$\begin{array}{l} \max_x \quad -\frac{1}{\epsilon} \log \det(\frac{1}{\epsilon} x^{-1}) + \eta \frac{1}{\epsilon} \langle -c, x \rangle \\ \text{s.t.} \quad A(x) = b \\ x \succ_{\text{star}} 0 \end{array}$$

$$\Rightarrow \begin{array}{l} \min_x \quad \langle c, x \rangle - \frac{1}{\epsilon} \log \det(x) \\ \text{s.t.} \quad A(x) = b \\ x \succ_{\text{star}} 0 \end{array} \quad \left. \vphantom{\begin{array}{l} \min_x \\ \text{s.t.} \end{array}} \right\} \text{primal barrier}$$

primal barrier and dual barrier have the same optimality condition

if $x(\epsilon), \lambda(\epsilon), v(\epsilon)$ solve the optimality condition,

then $x(\epsilon)$ solves the primal barrier

$\lambda(\epsilon), v(\epsilon)$ solves the dual barrier

KKT for dual barrier:

$$\text{primal feasible: } \lambda \succ_{\text{star}} 0 \quad A^T(v) + \lambda = c$$

$$\text{dual feasible: } x \succ_{\text{star}} 0 \quad A(x) = b$$

$$\text{gradient vanishes: } \lambda^{-1} = \frac{1}{\epsilon} x$$

Path-following Interior Point Method

generate $(x^k, \lambda^k, v^k) \approx (x(t^k), \lambda(t^k), v(t^k))$

key details $\left\{ \begin{array}{l} \text{proximity to central path} \\ \text{update: newton-like step} \end{array} \right.$

Local Neighborhoods of the Central Path

let $\mathcal{F}^0 = \{ (x, \lambda, v) : A(x) = b, A^T(v) + \lambda = c, x \succ 0, \lambda \succ 0 \}$ be the strictly feasible set

given $x \succ 0, \lambda \succ 0$ $u(x, \lambda) = \langle x, \lambda \rangle / n$

if x is primal feasible and λ, v is dual feasible $\langle c, x \rangle - \langle b, v \rangle = \langle c, x \rangle - \langle A(x), v \rangle = \langle x, \lambda \rangle$

$$d_{\mathcal{F}}(x, \lambda) = \| \lambda(x\lambda) - u(x, \lambda) \cdot \mathbf{1} \|_2$$

$$= \| x^{\lambda} x - u(x, \lambda) \cdot \mathbf{1} \|_2$$

$$= \| x^{\lambda} x^{\lambda} - u(x, \lambda) \cdot \mathbf{1} \|_2$$

(measure of proximity of $x\lambda$ and $u(x, \lambda) \cdot \mathbf{1}$)

$$\mathcal{N}_{\epsilon}(b) = \{ (x, \lambda, v) \in \mathcal{F}^0 : d_{\mathcal{F}}(x, \lambda) \leq \epsilon u(x, \lambda) \}$$

Newton Step

$$\begin{bmatrix} A^T(v) + \lambda - c \\ Ax - b \\ X\lambda \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ \frac{1}{2}I \end{bmatrix} \quad X, \lambda > 0$$

↓ natural newton step

$$\begin{bmatrix} 0 & A^T & I \\ A & 0 & 0 \\ \lambda & 0 & X \end{bmatrix} \begin{bmatrix} \Delta X \\ \Delta v \\ \Delta \lambda \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ \frac{1}{2}I - X\lambda \end{bmatrix}$$

ΔX and $\Delta \lambda$ may not be symmetric
($\frac{1}{2}I - X\lambda$ may not be symmetric)

Search Direction

$$r_z(x, \lambda, v) = \begin{bmatrix} A^T(v) + \lambda - c \\ Ax - b \\ X\lambda - \frac{1}{2}I \end{bmatrix} = \begin{bmatrix} \langle \bar{v}, A^T + \lambda - c \rangle \\ \langle A, x \rangle - b \\ \langle X\lambda, X\lambda \rangle - \frac{1}{2}I \end{bmatrix} \begin{matrix} r_d \text{ dual residual} \\ r_p \text{ primal residual} \\ r_c \text{ centrality residual} \end{matrix}$$

$$r_z(x + \Delta x, \lambda + \Delta \lambda, v + \Delta v) \approx r_z(x, \lambda, v) + D r_z \cdot \begin{bmatrix} \Delta x \\ \Delta \lambda \\ \Delta v \end{bmatrix} \Rightarrow$$

$$\begin{matrix} \textcircled{1} \\ \textcircled{2} \\ \textcircled{3} \end{matrix} \begin{bmatrix} 0 & A^T(v) & I \\ A(x) & 0 & 0 \\ \mathcal{E}(x) & 0 & f(x) \end{bmatrix} \begin{bmatrix} \Delta x \\ \Delta v \\ \Delta \lambda \end{bmatrix} = -r_z(x, \lambda, v) = - \begin{bmatrix} r_d \\ r_p \\ r_c \end{bmatrix}$$

$$\Delta x = \Sigma^{-1} (-r_c - f(\Delta \lambda)) \quad \text{from (3)}$$

$$= \Sigma^{-1} \left[-r_c - f(-rd - A^T(\Delta v)) \right] \quad \text{from (1)}$$

$$A \left(\Sigma^{-1} \left[-r_c - f(-rd - A^T(\Delta v)) \right] \right) = -r_p \quad \text{from (2)}$$

$$A \Sigma^{-1} (-r_c - f(-rd)) + A \Sigma^{-1} f A^T(\Delta v) = -r_p$$

$$\Delta v = (A \Sigma^{-1} f A^T)^{-1} \left[-r_p + A \Sigma^{-1} (r_c - f(-rd)) \right] \quad \text{solve for } \Delta v$$

$$\Delta \lambda = -rd - A^T(\Delta v) \quad \text{solve for } \Delta \lambda$$

$$\Delta x = \Sigma^{-1} (-r_c - f(\Delta \lambda)) \quad \text{solve for } \Delta x$$

Path Following Method

$$\begin{array}{l} \min. \langle C, x \rangle \\ \text{s.t. } Ax = b \\ x \geq 0 \end{array} \quad \longrightarrow \quad \begin{array}{l} \min t \langle C, x \rangle - \log \det(x) \\ \text{s.t. } Ax = b \end{array}$$

$$\tilde{J}(\tilde{x}, \tilde{v}) = t \langle C, \tilde{x} \rangle - \log \det(\tilde{x}) - \langle \tilde{v}, A\tilde{x} - b \rangle$$

$$\nabla_{\tilde{x}} \tilde{J} = tC - \tilde{x}^{-1} - A^T \tilde{v} = 0$$

$$C - \underbrace{\tilde{x}^{-1}}_{\frac{1}{t} \tilde{x}^*} - A^T \tilde{v} = 0 \quad \Rightarrow \quad \nabla_{\tilde{x}} \tilde{J} = \frac{d}{dt} \langle C, x \rangle - \langle \lambda, x \rangle - \langle \tilde{v}, A(x) - b \rangle$$

if \tilde{x}, \tilde{v} solves the log-barrier, then \tilde{x} minimize the Lagrangian

$$J(x, \lambda, v) = \langle C, x \rangle - \langle \lambda, x \rangle - \langle v, A(x) - b \rangle$$

$$\text{where } \lambda = \frac{1}{t} \tilde{x}^{-1}, \quad v = \tilde{v} / t$$

$$\begin{aligned} p^* &\geq g(\lambda, v) = \inf_x J(x, \lambda, v) \\ &= \langle C, \tilde{x} \rangle - \langle \lambda, \tilde{x} \rangle - \langle v, A(\tilde{x}) - b \rangle \\ &= \langle C, \tilde{x} \rangle - \frac{1}{t} \langle \tilde{x}^{-1}, \tilde{x} \rangle \\ &= \langle C, \tilde{x} \rangle - \frac{n}{t} \quad \text{duality gap} \end{aligned}$$

where x is primal feasible, λ, v is dual feasible

$$\langle \lambda, x \rangle = \frac{n}{t} \quad \text{is the duality gap}$$

$$t = \frac{n}{\langle \lambda, x \rangle}$$

while the:

$$\text{if } A^T v + \lambda - C = 0 \quad \text{and} \quad Ax = b \quad \text{and} \quad \langle \lambda, x \rangle \leq \epsilon$$

return x, λ, v

$$t = \frac{n}{\langle \lambda, x \rangle} \cdot \alpha \quad // \quad \alpha \text{ controls the advancement of central path}$$

solve the system (Solve the log-barrier)

$$\begin{cases} A^T(v + \alpha v) + (\lambda + \alpha \lambda) - C = 0 \\ A(x + \alpha x) = b \\ (x + \alpha x)(\lambda + \alpha \lambda) = \frac{1}{t} I \end{cases}$$

line search to get s

$$\begin{aligned} x &+ = s \cdot \alpha x \\ \lambda &+ = s \cdot \alpha \lambda \\ v &+ = s \cdot \alpha v \end{aligned}$$